

Interval Trees

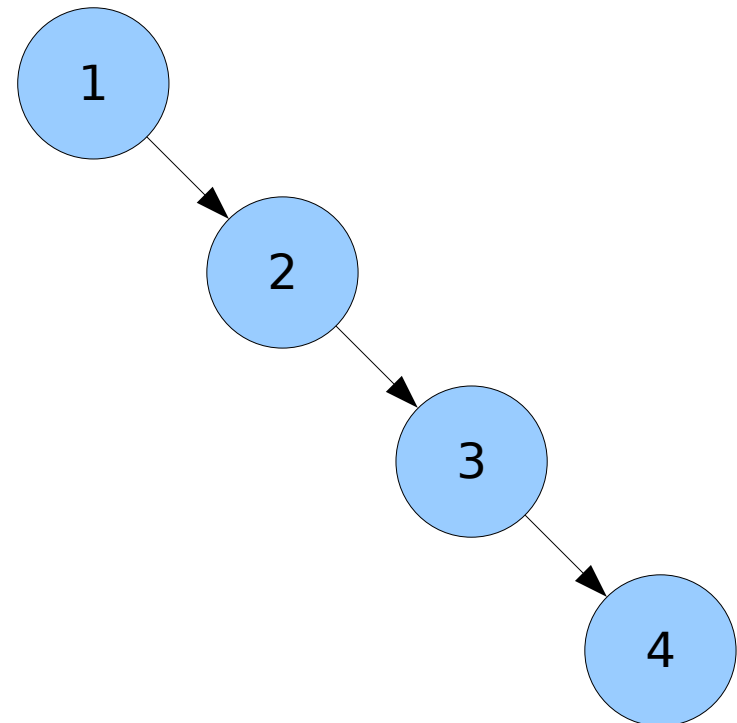
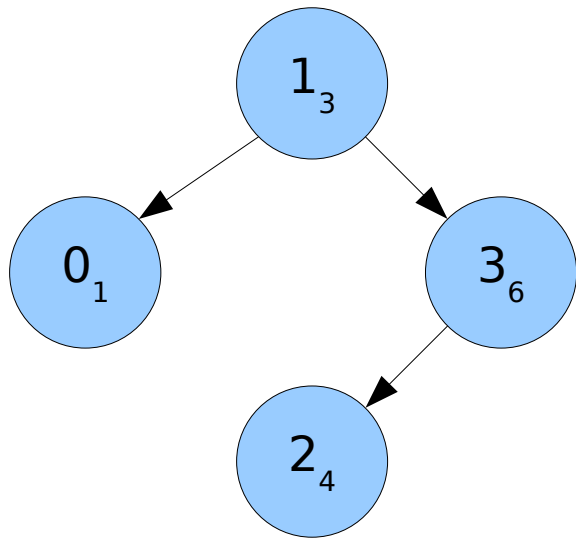
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Problem

- Given a collection of items i , each with value V_i
- Want to answer many queries of the form:
How many items are in the interval $[x, y]$?

Binary Tree

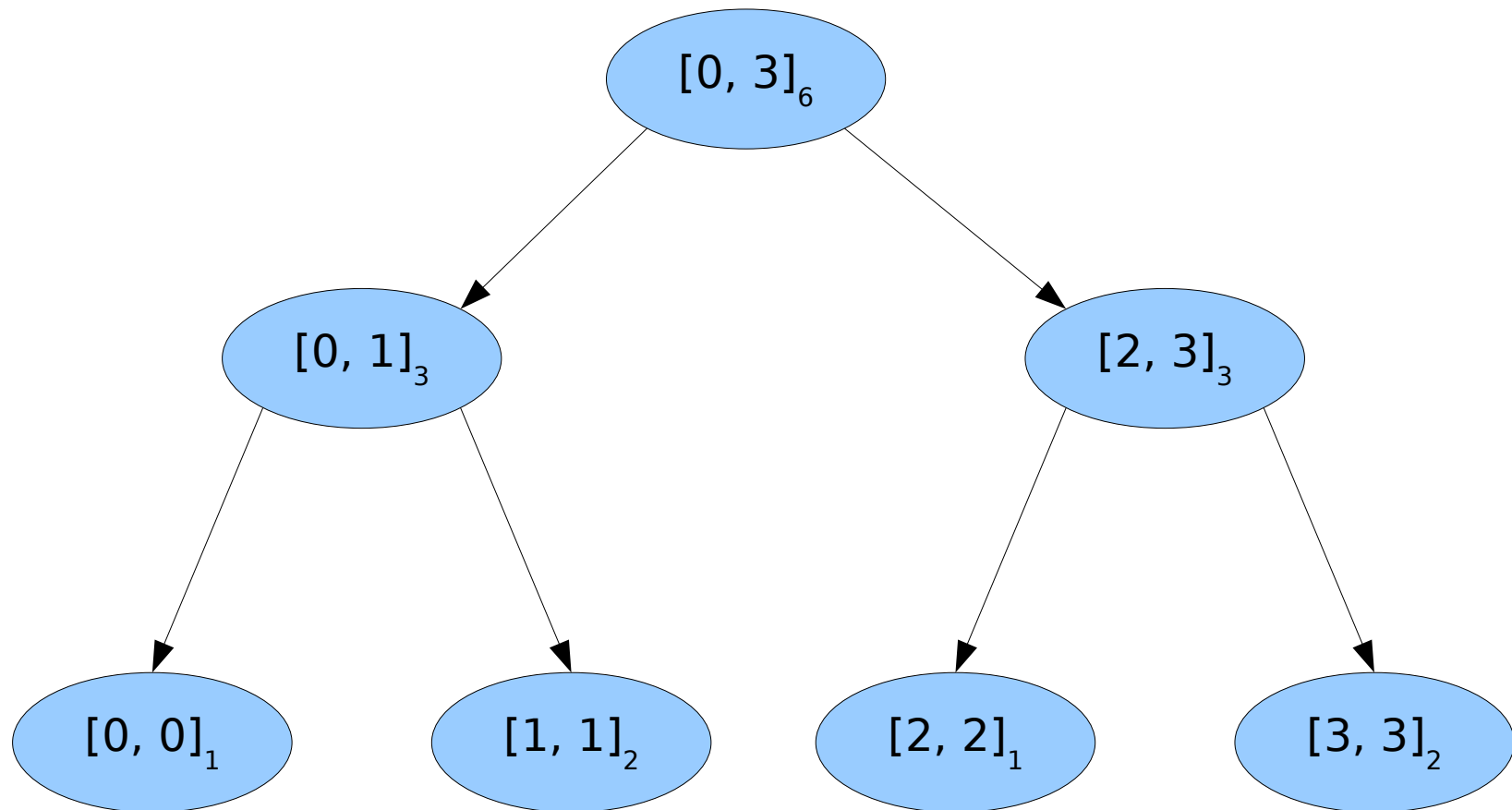
- At each node V_i , store the count o of interval $[0, V_i]$
- Problem: unbalanced trees:



Radix Tree

- Root node contains the interval $[0, \text{max}]$
- A node containing the interval $[A, B]$ has children with interval $[A, (A+B)/2]$ and $[(A+B)/2 + 1, B]$
- Each node holds the count over its range
- Every number has a well-defined position

Radix Tree: Example



Radix Tree: Update

- To insert the value V_i :
 - Increment the root node's count
 - If $V_i \leq (A+B)/2$
 - Insert V_i into the left child
 - Else
 - Insert V_i into the right child
- Similar process for deletion

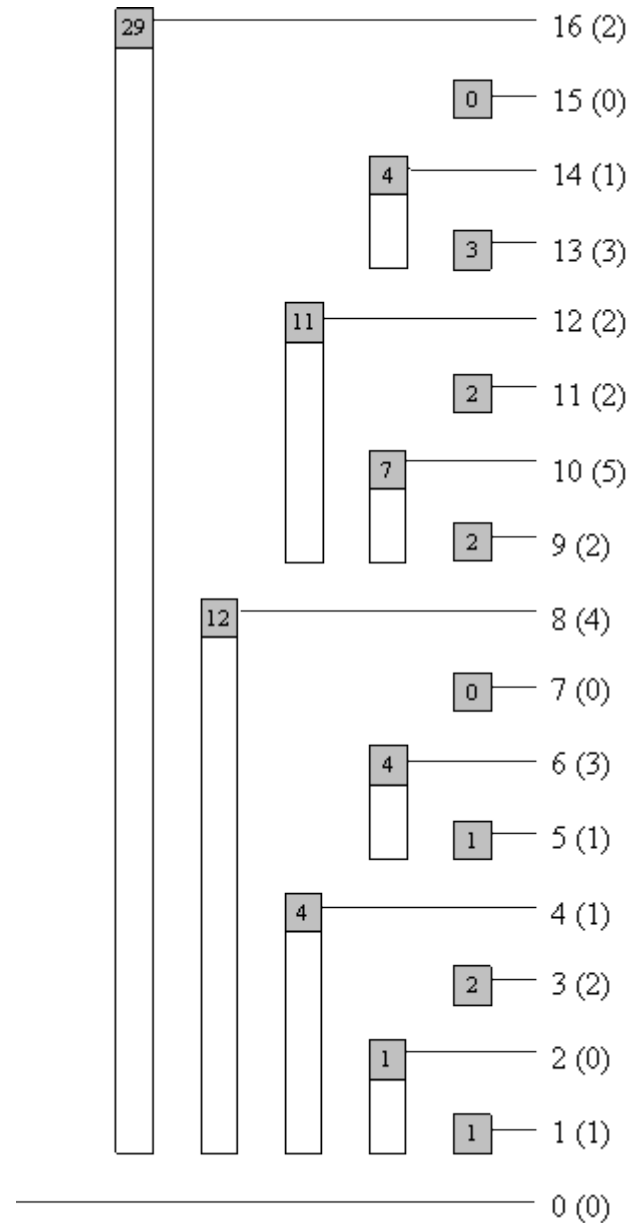
Radix Tree: Query

- To find the number of items in the interval $[x, y]$:
 - If $[A, B]$ covers $[x, y]$
 - return root node's count
 - Let $sum = 0$
 - If $[A, (A+B)/2]$ overlaps $[x, y]$
 - $sum +=$ count in left child over $[x, y]$
 - If $[(A+B)/2 + 1, B]$ overlaps $[x, y]$
 - $sum +=$ count in right child over $[x, y]$
 - Return sum

Binary Indexed Trees

- Any number can be represented as the sum of powers of two
- Assume intervals are of the form $[1, x]$
- An interval count can be represented in a similar manner
- $[1, 13] = [1, 8] + [9, 12] + [13, 13]$
- So we only store counts of interval of the form $[x - 2^r + 1, x]$

Binary Indexed Trees: Example



Binary Indexed Trees

- $13 = 1101_2$
- $c[1101] = \text{tree}[1101] + \text{tree}[1100] + \text{tree}[1000]$
- To isolate last 1 in binary form: $n \& -n$
 - Proof in reference article

Binary Indexed Trees: Update

```
void update(int idx ,int val) {  
    while (idx <= MaxVal) {  
        tree[idx] += val;  
        idx += (idx & -idx);  
    }  
}
```

Binary Indexed Trees: Query

```
int read(int idx) {  
    int sum = 0;  
    while (idx > 0) {  
        sum += tree[idx];  
        idx -= (idx & -idx);  
    }  
    return sum;  
}
```

Sample Problem: TopCoder SRM 310

- n cards face down on table
 - $T\ i\ j$: turn cards from index i to index j , include i -th and j -th card
 - $Q\ i$: is the i -th card face?)

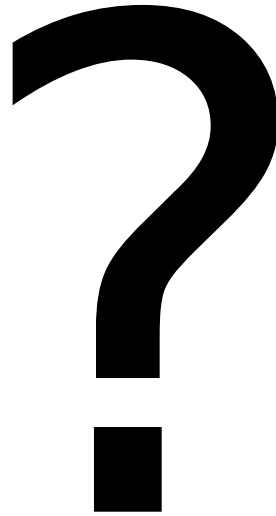
Sample Problem: Solution

- Array f (of length $n + 1$) holds our BIT
- $f[i]++$ and $f[j + 1]--$
- For each card k between i and j , sum $f[1] + f[2] + \dots + f[k]$ will be incremented
- For each query, the answer is $f[i] \% 2$

Analysis

- Binary Tree: common, but balancing can be difficult and expensive
- Radix Tree: guaranteed $O(\log N)$ for update and query
- Binary Index Tree: Same as Radix Tree, but very short code

Questions



References

- <http://www.topcoder.com/tc?module=Static&d1=tutorials&d2=binaryIndexedTrees>