

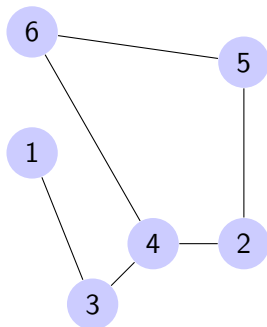
Intro to Graph Theory

2014 IOI Camp 1

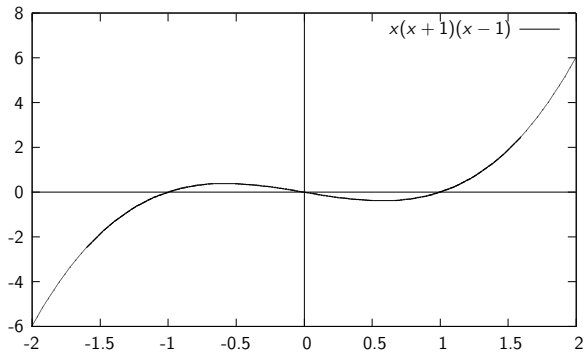
Robert Spencer

December 11, 2013

This is a graph:



This is not a graph:



Definition

A graph is a collection of *nodes* connected by *edges* which may or may not be *directed* and/or *weighted*

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Examples of graphs:

- A computer network (non-directed, non-weighted)
- A road map (non-directed, weighted)
- Winners in a chess tournament (directed, non-weighted)
- Payments in an economy (weighted, directed)

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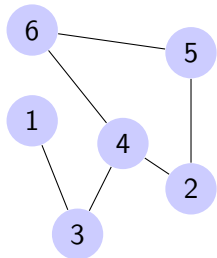
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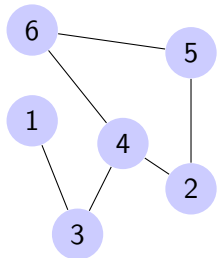
Can you find a path and a cycle?



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Example Answer:

Path: 1-3-4-2

Cycle: 4-2-5-6-4

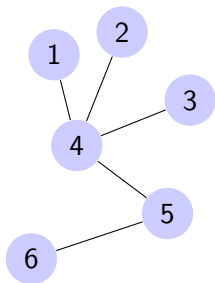
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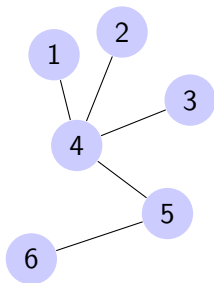
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Example



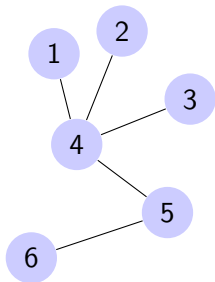
Theorem

A tree of n vertices has $n - 1$ edges.

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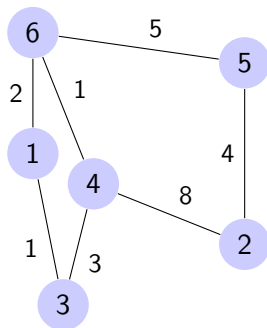
Theorem

A tree of n vertices has $n - 1$ edges.

Proof.

Induction. Start with one vertex, and add subsequent ones. □

Weights are placed on edges, and can represent anything (lengths, costs, etc.)



Graph Representations

How do we represent a graph?

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Lists of Neighbours

$[(3,1), (6,2)]$

$[(4,8), (5,4)]$

$[(1,1), (4,3)]$

$[(2,8), (3,3), (6,1)]$

$[(2,4), (6,5)]$

$[(5,5)]$

Memory $O(E)$

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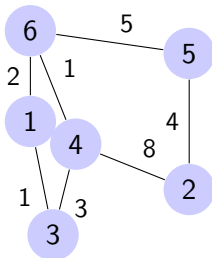
Adjacency Matrix

$$\begin{pmatrix} 0 & 0 & 1 & 0 & 0 & 2 \\ 0 & 0 & 0 & 8 & 4 & 0 \\ 1 & 0 & 0 & 3 & 0 & 0 \\ 0 & 8 & 3 & 0 & 0 & 1 \\ 0 & 4 & 0 & 0 & 0 & 5 \\ 0 & 0 & 0 & 1 & 5 & 0 \end{pmatrix}$$

Memory $O(N^2)$

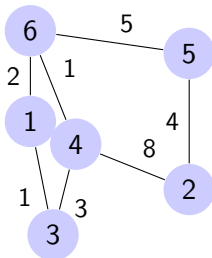
Traversal

Sometimes we want to visit all the nodes in a graph in a particular order. For example to search for a path/destination



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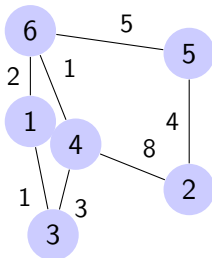
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We may visit nodes more than once, as there may be more than one path. E.g. to get from 1 to 2, we may visit 4 twice: 1-6-4-2 or 1-3-4-2.

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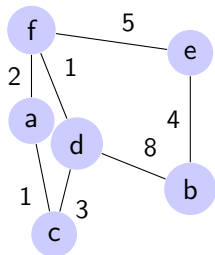
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Often this is used to find the shortest route between two or more nodes.

Depth First Search

Depth First Search (DFS) visits the nodes as far as it can before backtracking (without visiting nodes more than once).

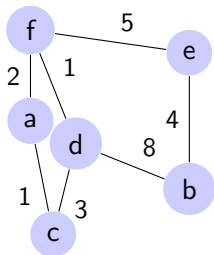
Sample Graph:



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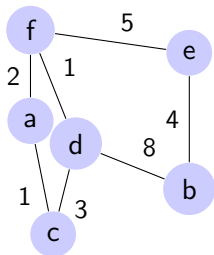
Pseudocode:

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def DFS(currNode, finalNode)
  if currNode==finalNode then
    return success
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  foreach neighbour of currNode do
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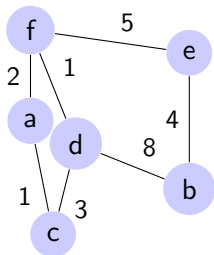
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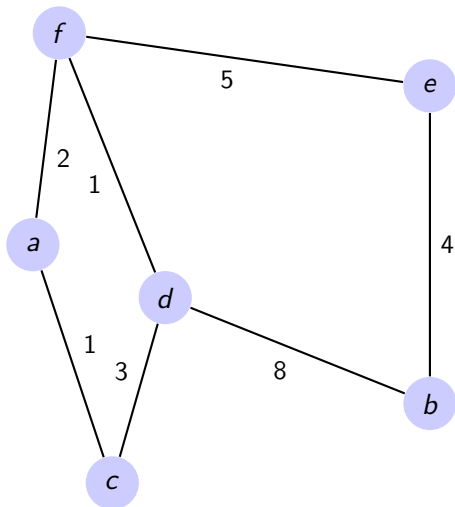


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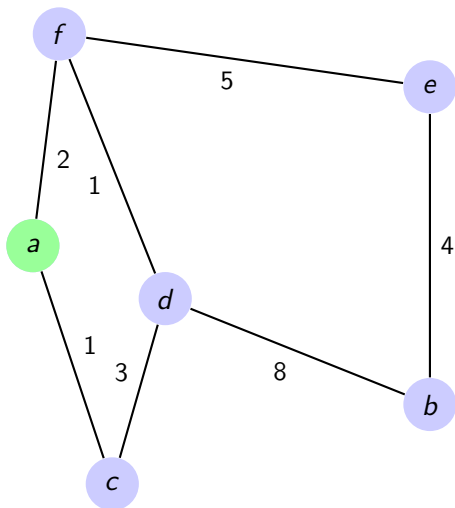
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a-c-d-b-e-f-f-e-b-f-d-c-b-e

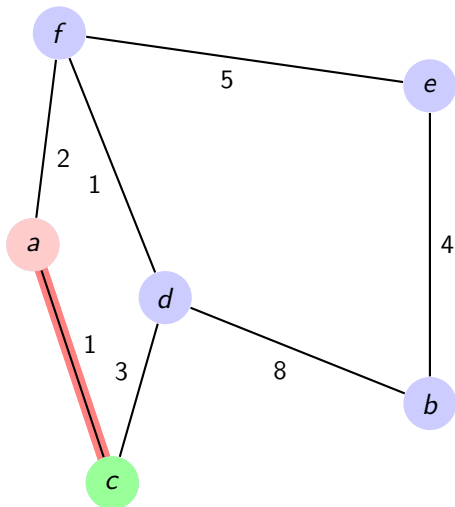
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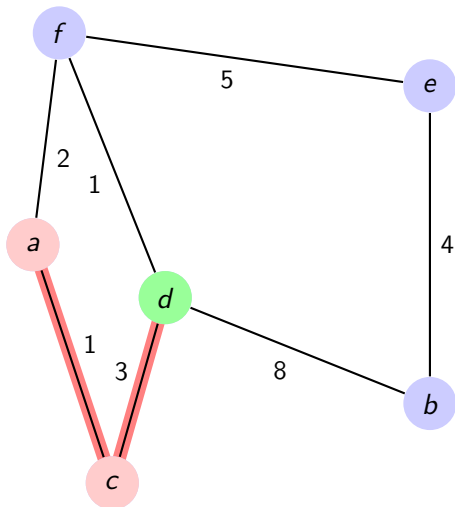
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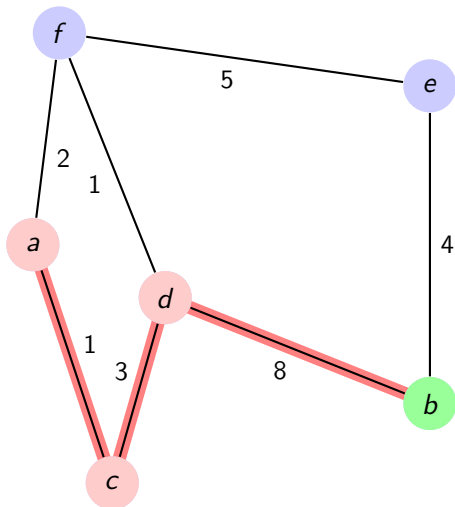
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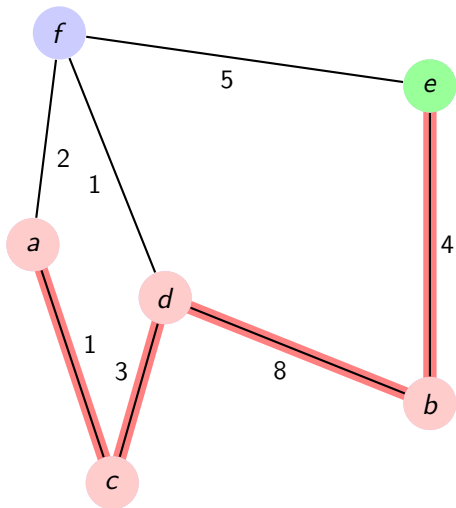
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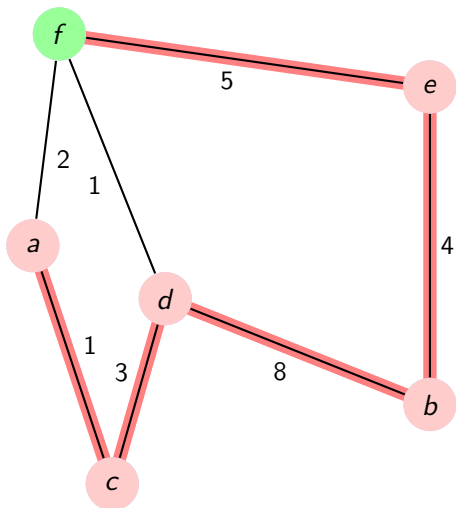
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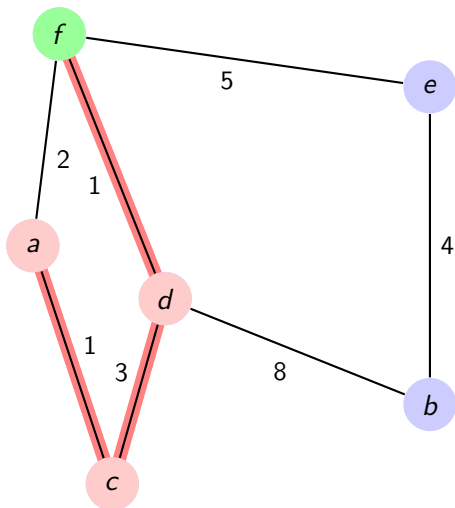
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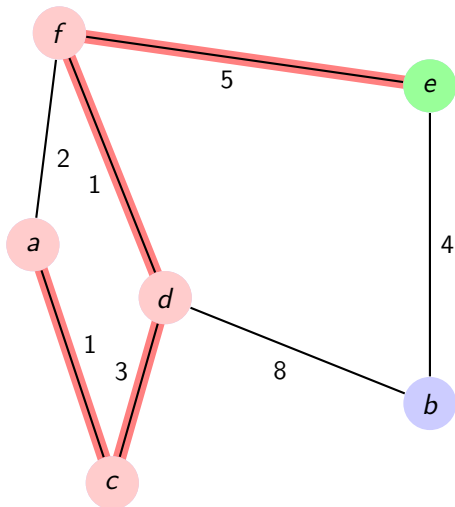
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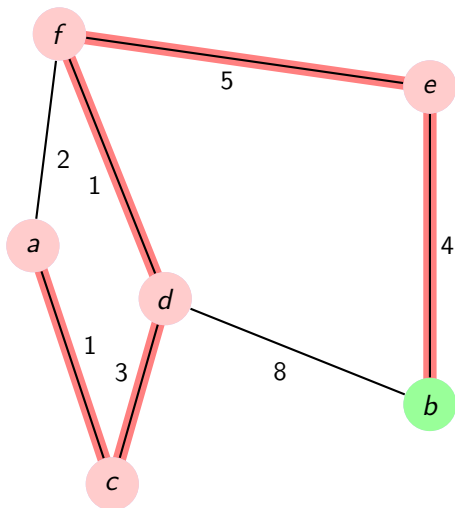
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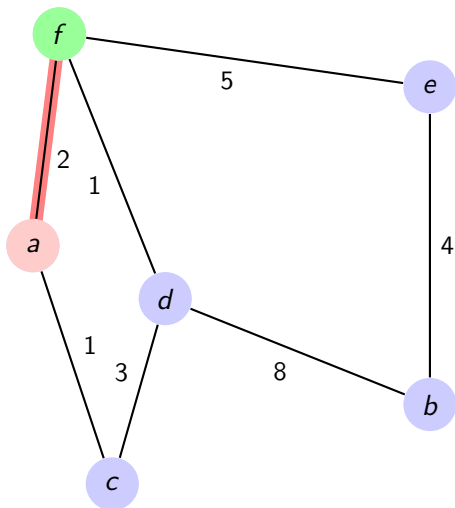
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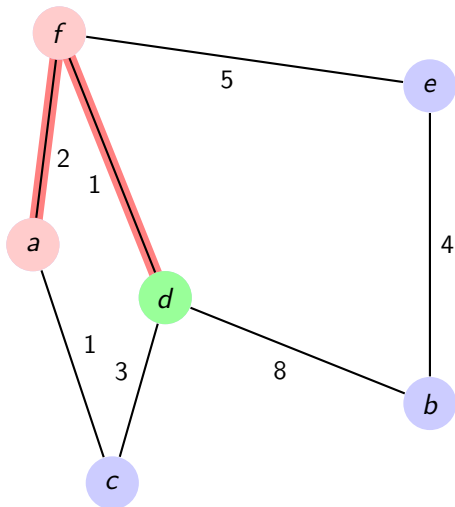
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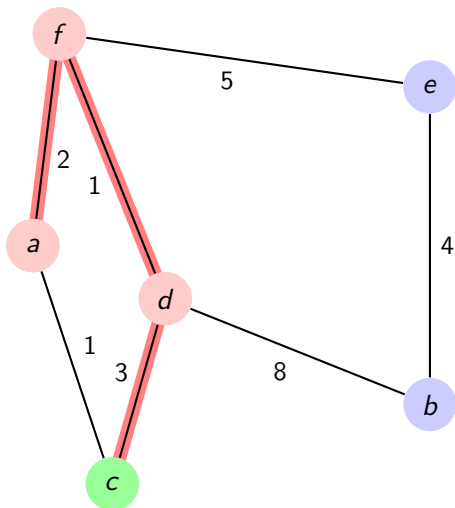
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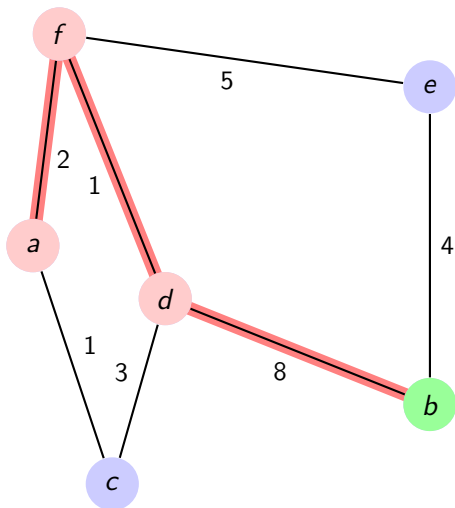
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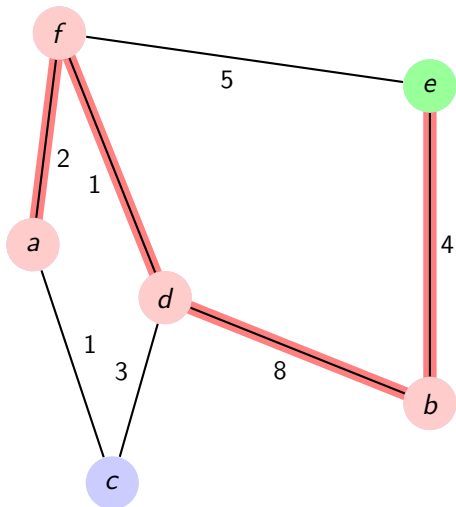
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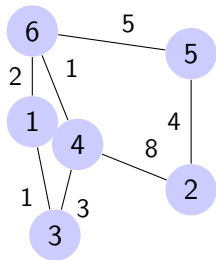
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Breadth First Search

Breadth First Search (DFS) visits the nodes “in parallel” without backtracking.

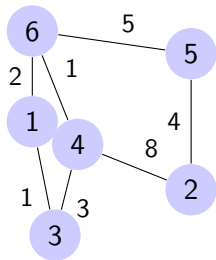
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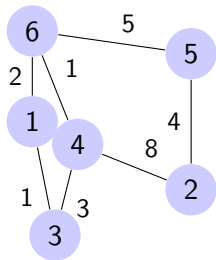
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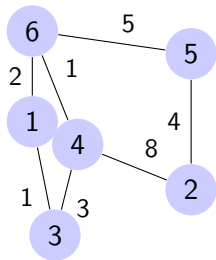
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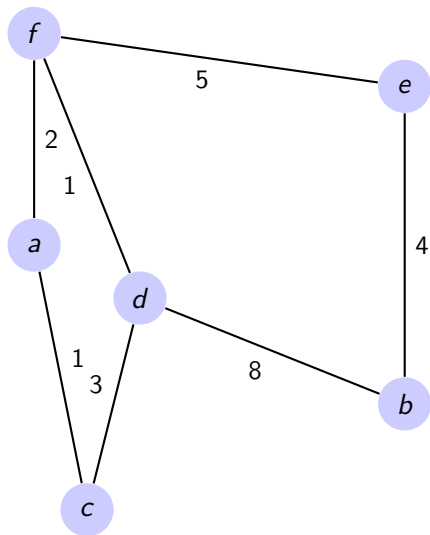


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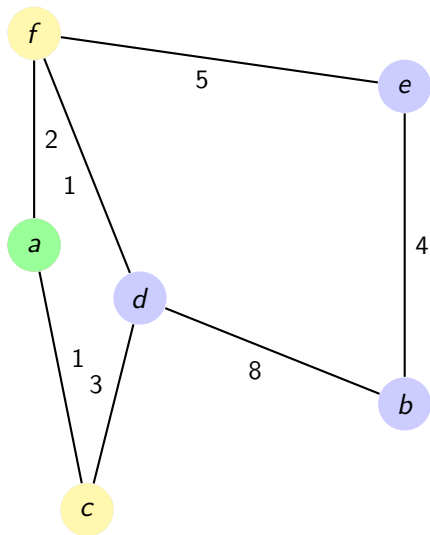
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Nodes (starting from 1) will be visited in this order: 1-3-6-4-5-2

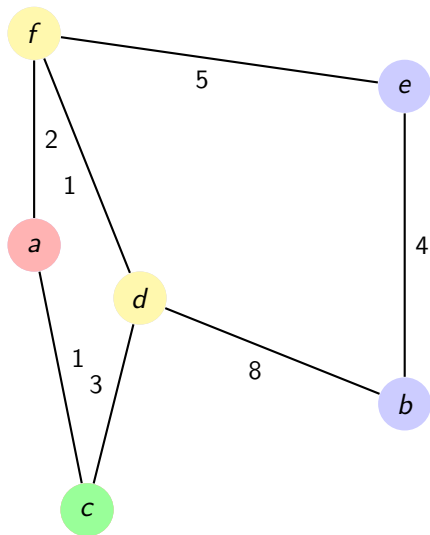
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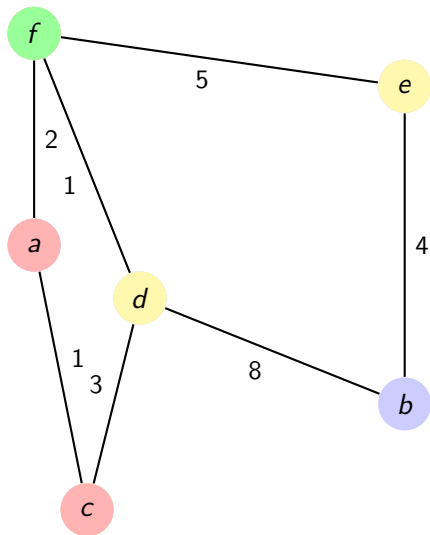
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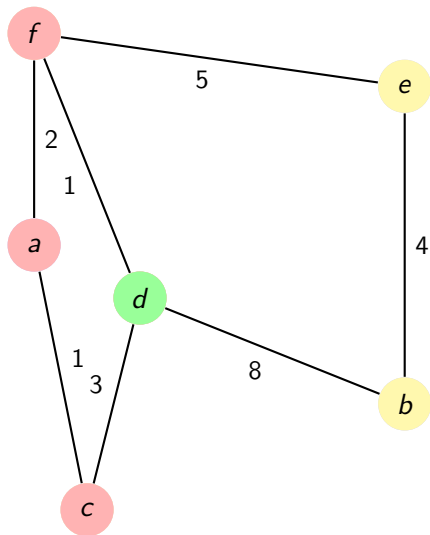
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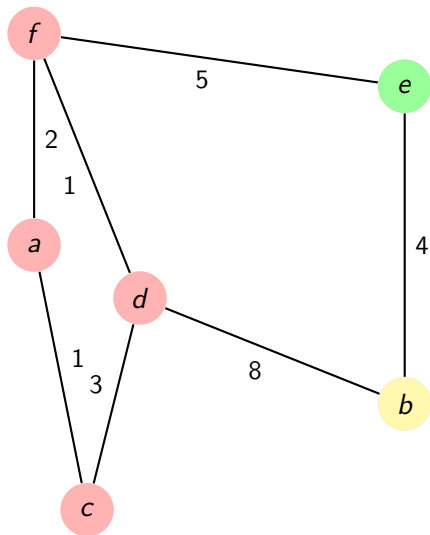
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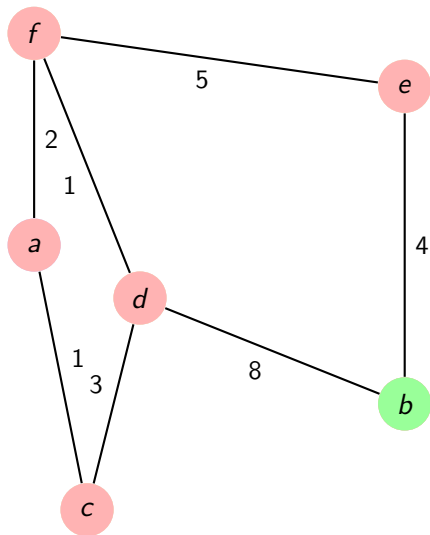
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Breadth First Search Example



Dijkstra's Algorithm

Dijkstra's Algorithm finds the shortest distance from one node to all others. It is basically a BFS with a priority queue.

Pseudocode:

```
set all distances INF
add (0, startNode) to queue
while queue not empty do
    currDists, currNode = queue.pop
    distances[currNode] = currDist
    for neighbour, distance in adjacent[currNode] do
        possNewDist = distances[currNode] + distance
        if distances[neighbour] > possNewDist then
            update neighbour to weight possNewDist in queue
```

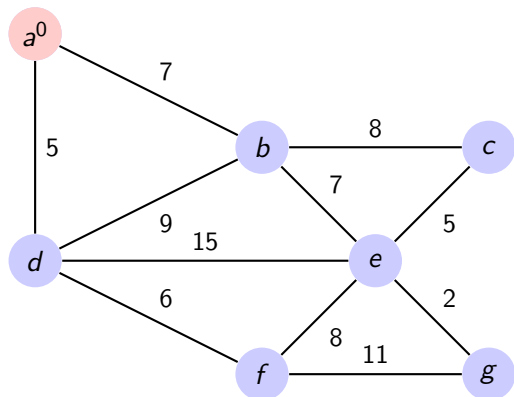
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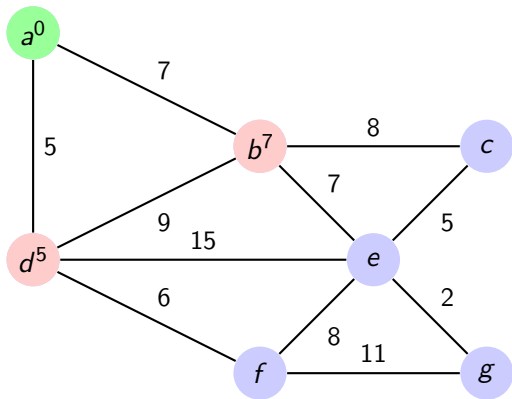
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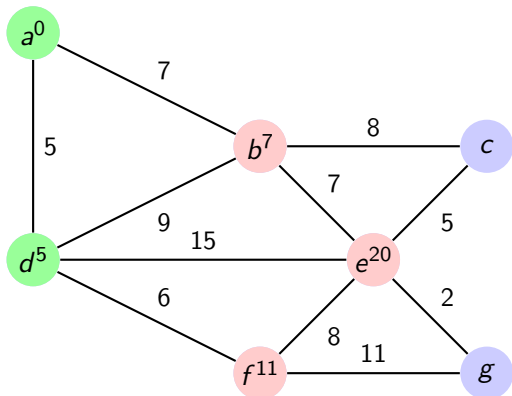
Queue: $\{(a, 0)\}$

Dijkstra Example



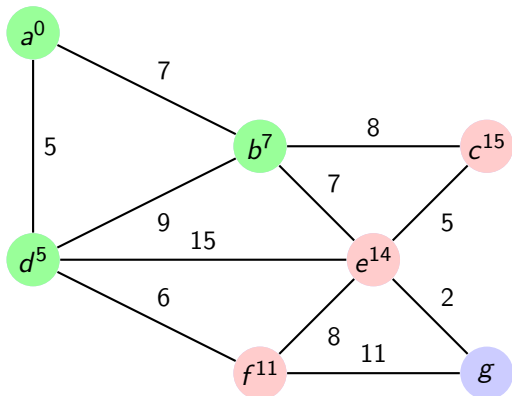
Queue: $\{(d, 5), (b, 7)\}$

Dijkstra Example



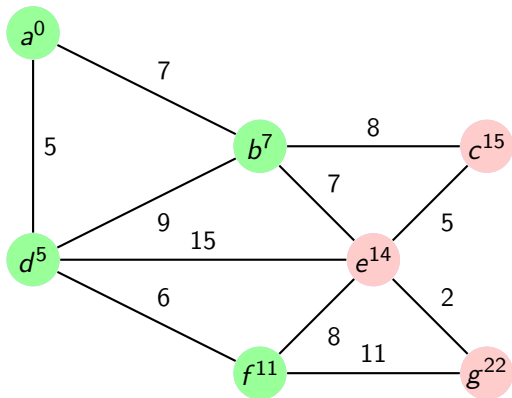
Queue: $\{(b, 7), (f, 11), (e, 20)\}$

Dijkstra Example



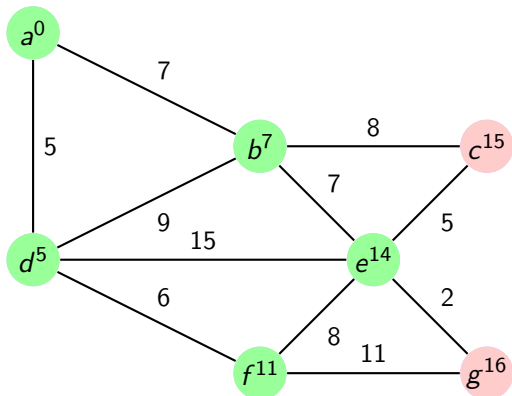
Queue: $\{(f, 11), (e, 14), (c, 15)\}$

Dijkstra Example



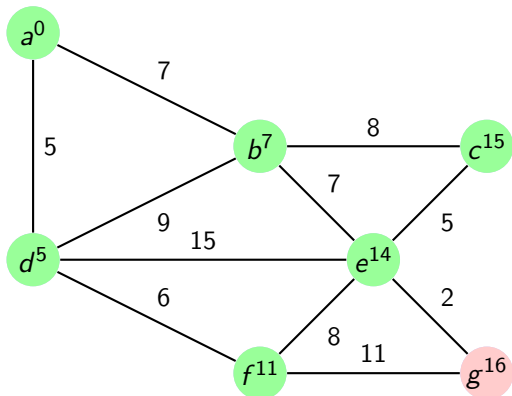
Queue: $\{(e, 14), (c, 15), (g, 22)\}$

Dijkstra Example



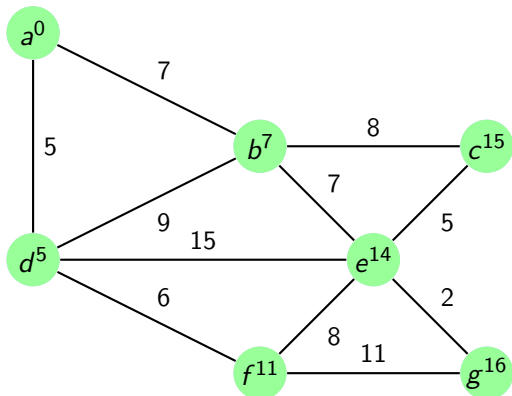
Queue: $\{(c, 15), (g, 16)\}$

Dijkstra Example



Queue: $\{(g, 16)\}$

Dijkstra Example

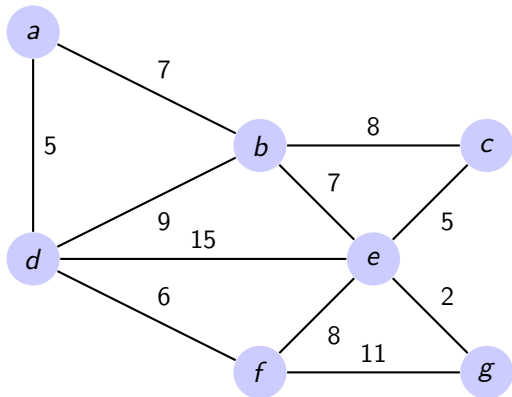


Queue: {}

Minimum Spanning Tree

Definition

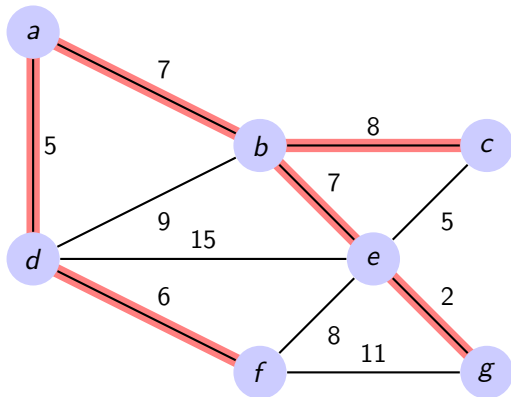
A *minimum spanning tree* is a subset of edges in a weighted undirected graph such that the edges form a tree containing all the nodes, and the sum of the weights of the tree is minimal.



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Prim's Algorithm

Prim's Algorithm finds the minimum spanning tree from a given graph.

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- Set all vertices to “not in the tree” except a starting vertex.
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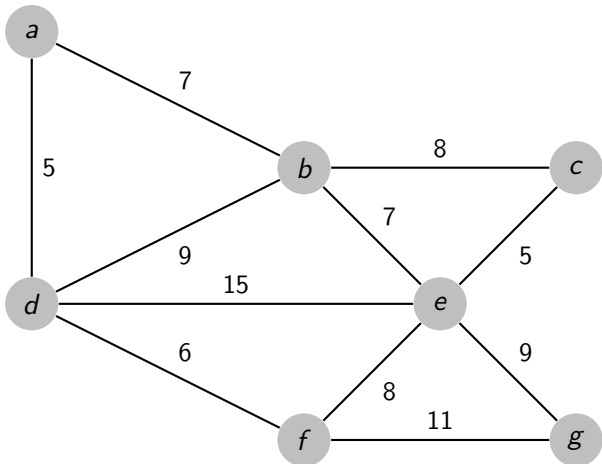
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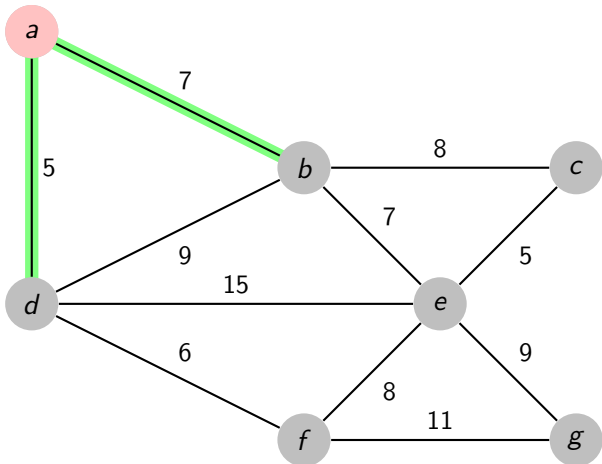
Technical notes

- Keep a priority queue of edges. Each step pull off an edge, check if it joins a new vertex. If it does, add all the edges from that vertex to the queue.
- Runs in $O(E \log V)$ with a binary heap as priority queue.

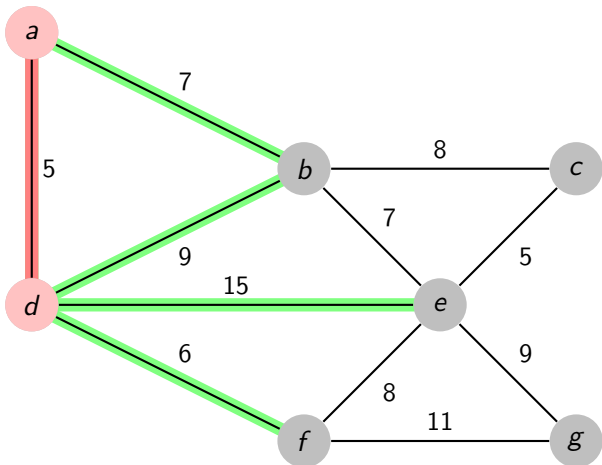
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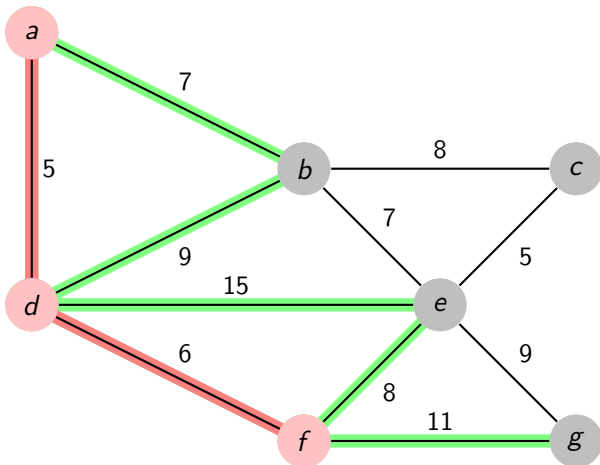
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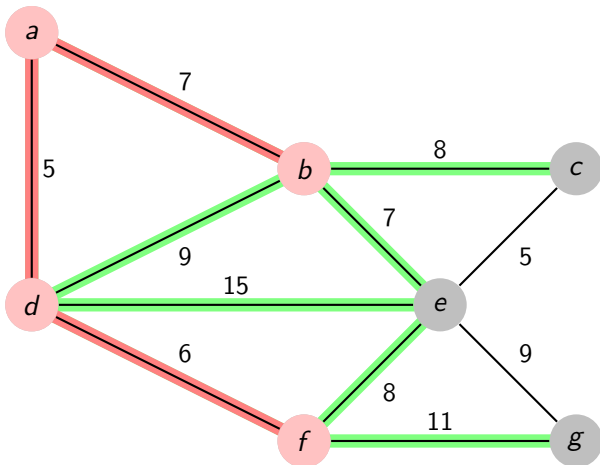
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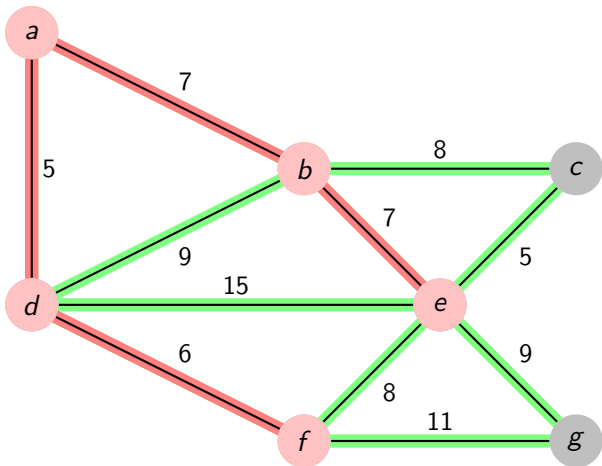
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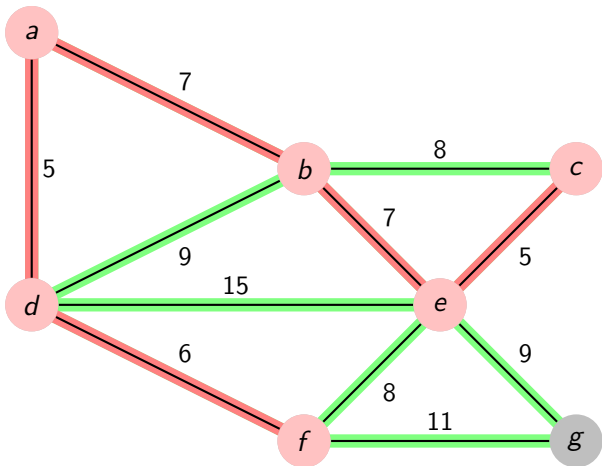
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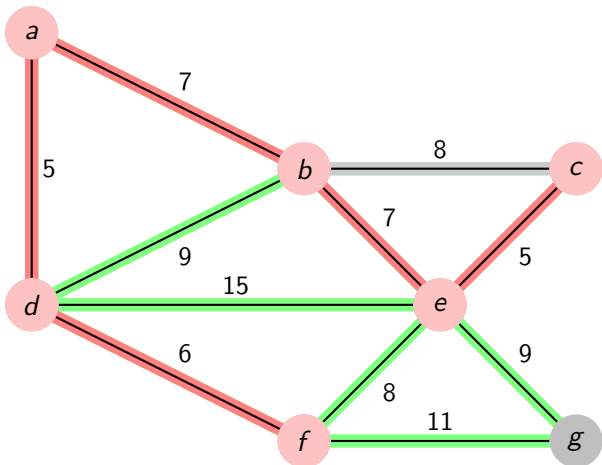
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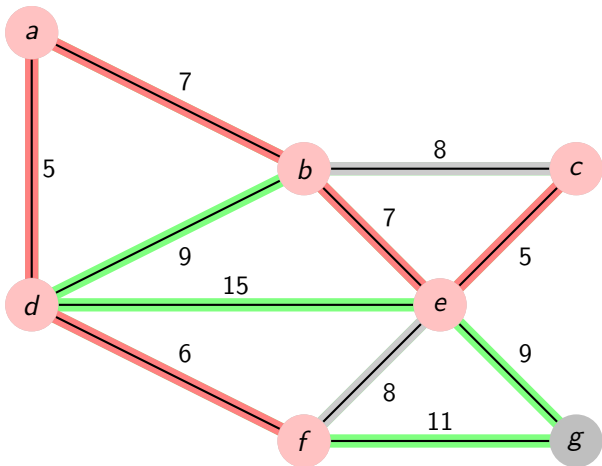
Prim's Algorithm Example



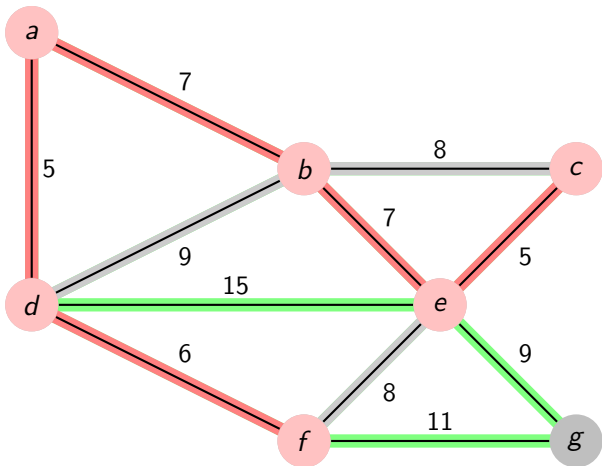
Prim's Algorithm Example



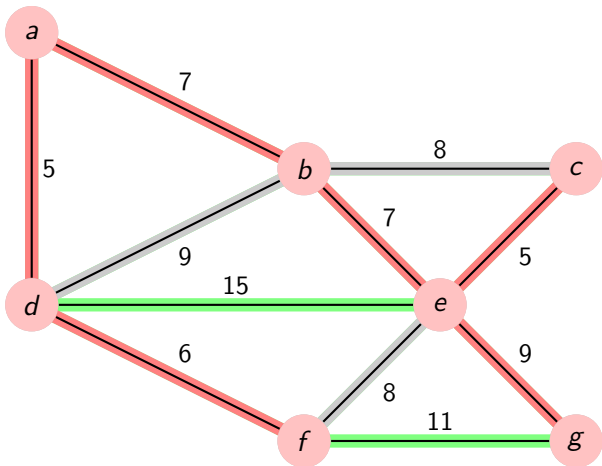
Prim's Algorithm Example



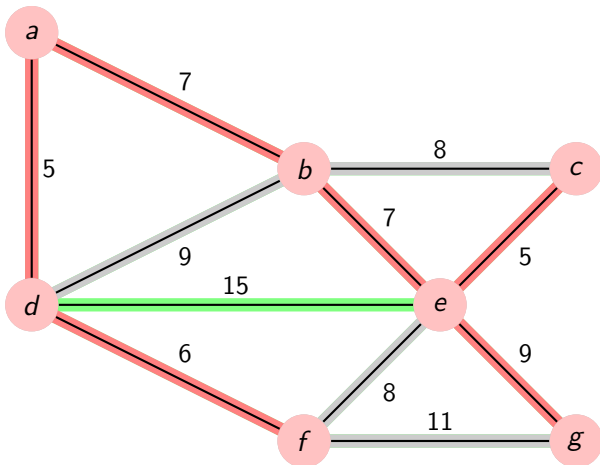
Prim's Algorithm Example



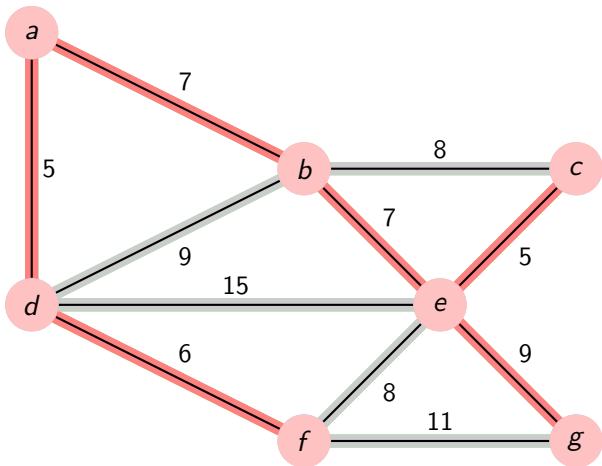
Prim's Algorithm Example



Prim's Algorithm Example



Prim's Algorithm Example



Kruskal's Algorithm

Kruskal's Algorithm is the “dual” of Prim's. It also finds the minimum spanning tree.

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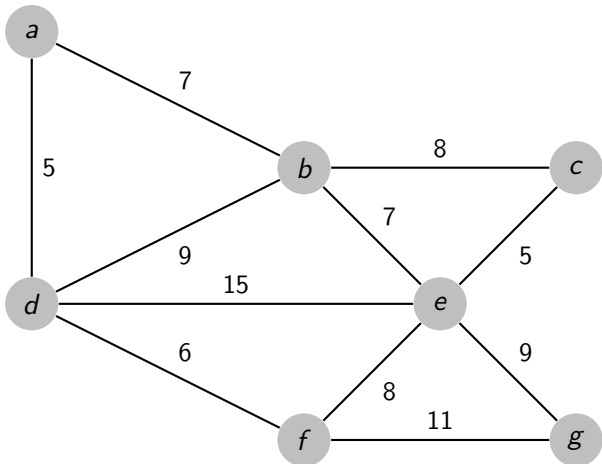
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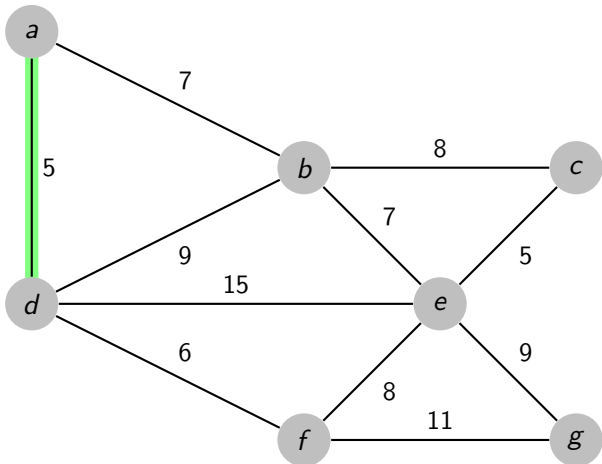
Technical Notes:

- Use “union-find” to hold the different trees.
- Complexity $O(E \log E)$

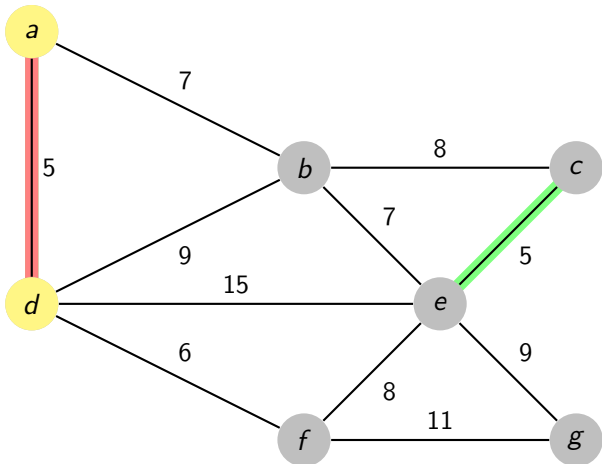
Kruskal's Algorithm Example



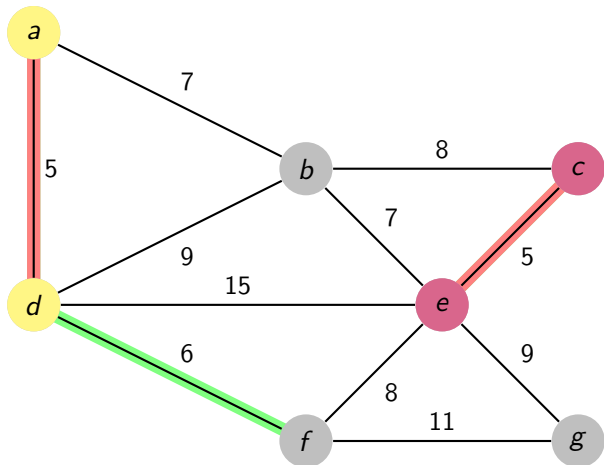
Kruskal's Algorithm Example



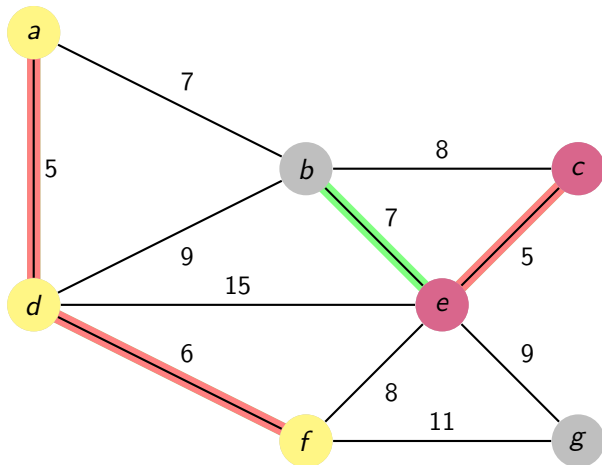
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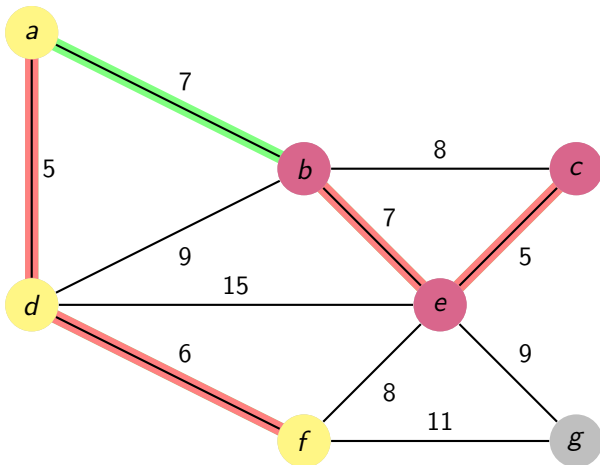
Kruskal's Algorithm Example



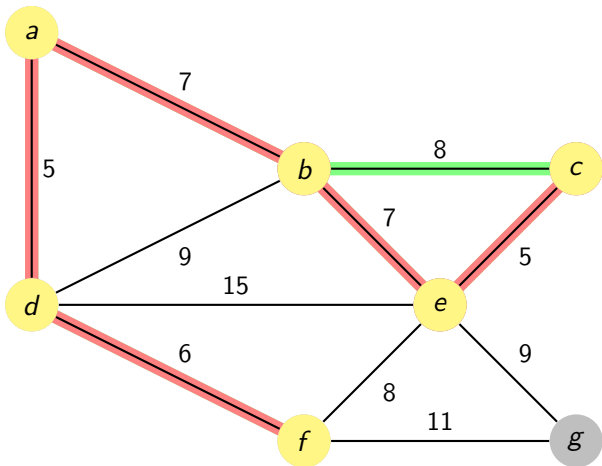
Kruskal's Algorithm Example



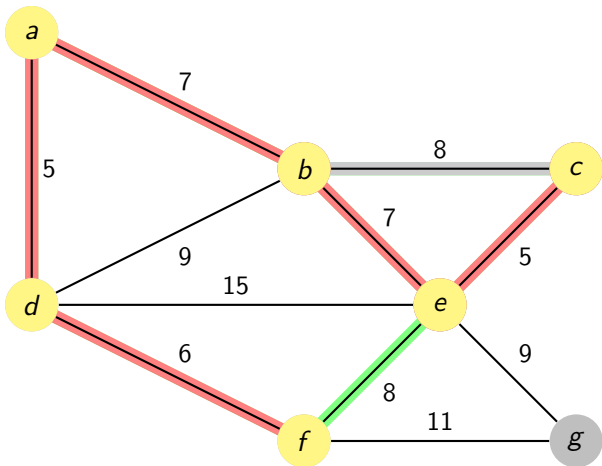
Kruskal's Algorithm Example



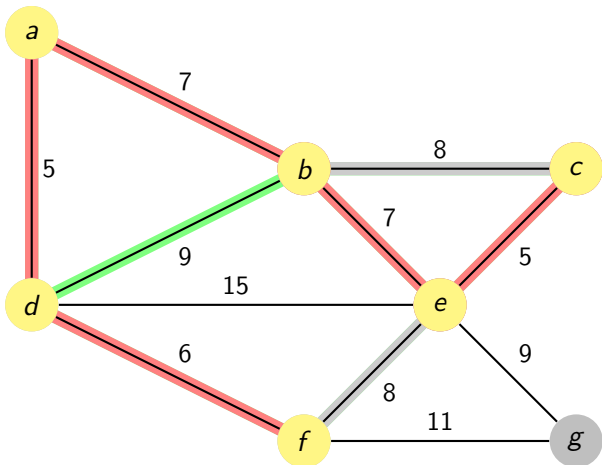
Kruskal's Algorithm Example



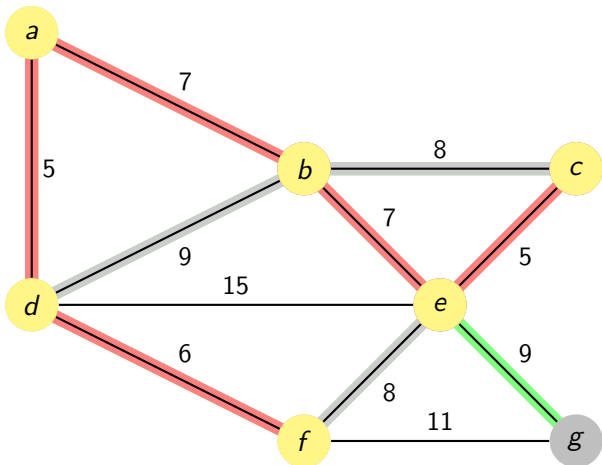
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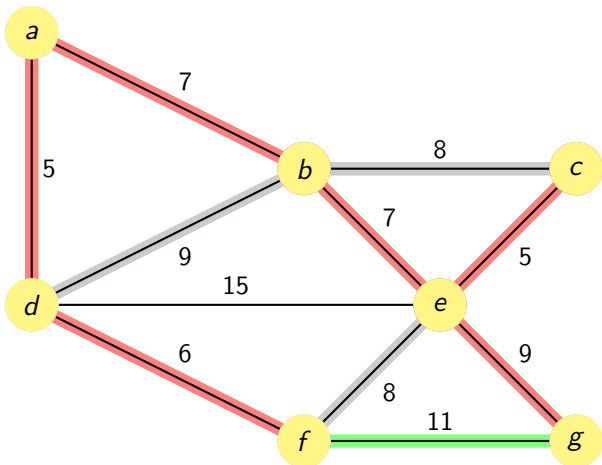
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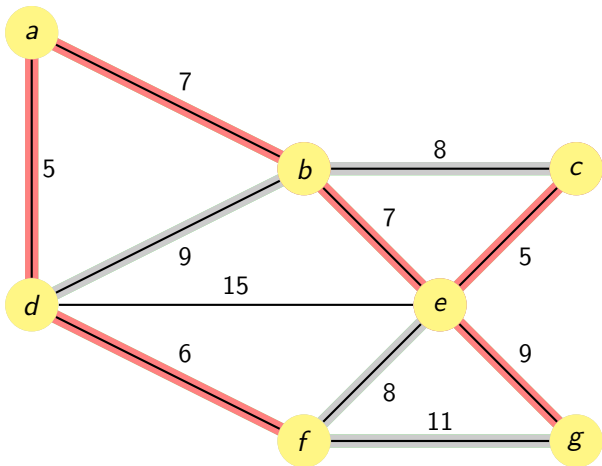
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